

Homework Assignment #9

Due: August 12, 2020 at 12:00 noon

1. In this question, we consider the problem of building a shared implementation of a stack. We will only be concerned with stacks that are limited to contain at most N items at any time. In other words, you can assume (as a precondition of PUSH) that no process will try to PUSH a value on to a stack that already contains N elements.

- (a) Consider the following implementation from read/write registers. It uses a shared array $A[1..N]$ of read/write registers and a shared read/write register Top . Assume the stack is initially empty and that Top is initialized to 0. Each process also uses a *local* variable t .

```

1  PUSH( $x$ )
2       $t \leftarrow Top.read$ 
3       $Top.write(t + 1)$ 
4       $A[t + 1].write(x)$ 
5      return ACK
6  end PUSH

7  POP
8       $t \leftarrow Top.read$ 
9      if  $t = 0$  then return EMPTY
10     else
11          $Top.write(t - 1)$ 
12         return  $A[t].read$ 
13     end if
14 end POP

```

Is this implementation correct if only one process accesses the stack at a time? Briefly justify your answer.

- (b) Now, suppose that the implementation in part (a) can be accessed by two processes concurrently. Consider an execution where each of the two processes perform a PUSH operation followed by a POP. Furthermore, suppose both PUSH operations execute line 2 before either operation executes line 3. Construct the rest of the execution so that the execution is *not* linearizable. Conclude that the implementation in part (a) is not linearizable.
- (c) Now suppose we discover that the hardware that we are using for our stack implementation also provides two machine instructions called `fetch&inc` and `fetch&dec`. A `fetch&inc` on a memory location allows us to atomically fetch the old value stored in the memory location and increment it by 1. In other words, a `fetch&inc` instruction changes the state of the memory location from the value v to $v + 1$ and returns v . Similarly, a `fetch&dec` instruction atomically fetches the old value stored in the memory location and decrements it by 1. (Assume that the value in the memory location is a non-negative integer, and that a `fetch&dec` on a memory location whose value is 0 returns 0 and does not change the value stored there, so that the value never becomes negative.) We use these new instructions to rewrite our implementation in part (a) as follows.

```

15 PUSH( $x$ )
16      $t \leftarrow Top.fetch\&inc$ 
17      $A[t + 1].write(x)$ 
18     return ACK
19 end PUSH

```

```
20 POP
21    $t \leftarrow Top.fetch\&dec$ 
22   if  $t = 0$  then return EMPTY
23   else return  $A[t].read$ 
24   end if
25 end POP
```

Is this a linearizable implementation of a stack (for two processes)? Show that your answer is correct.