

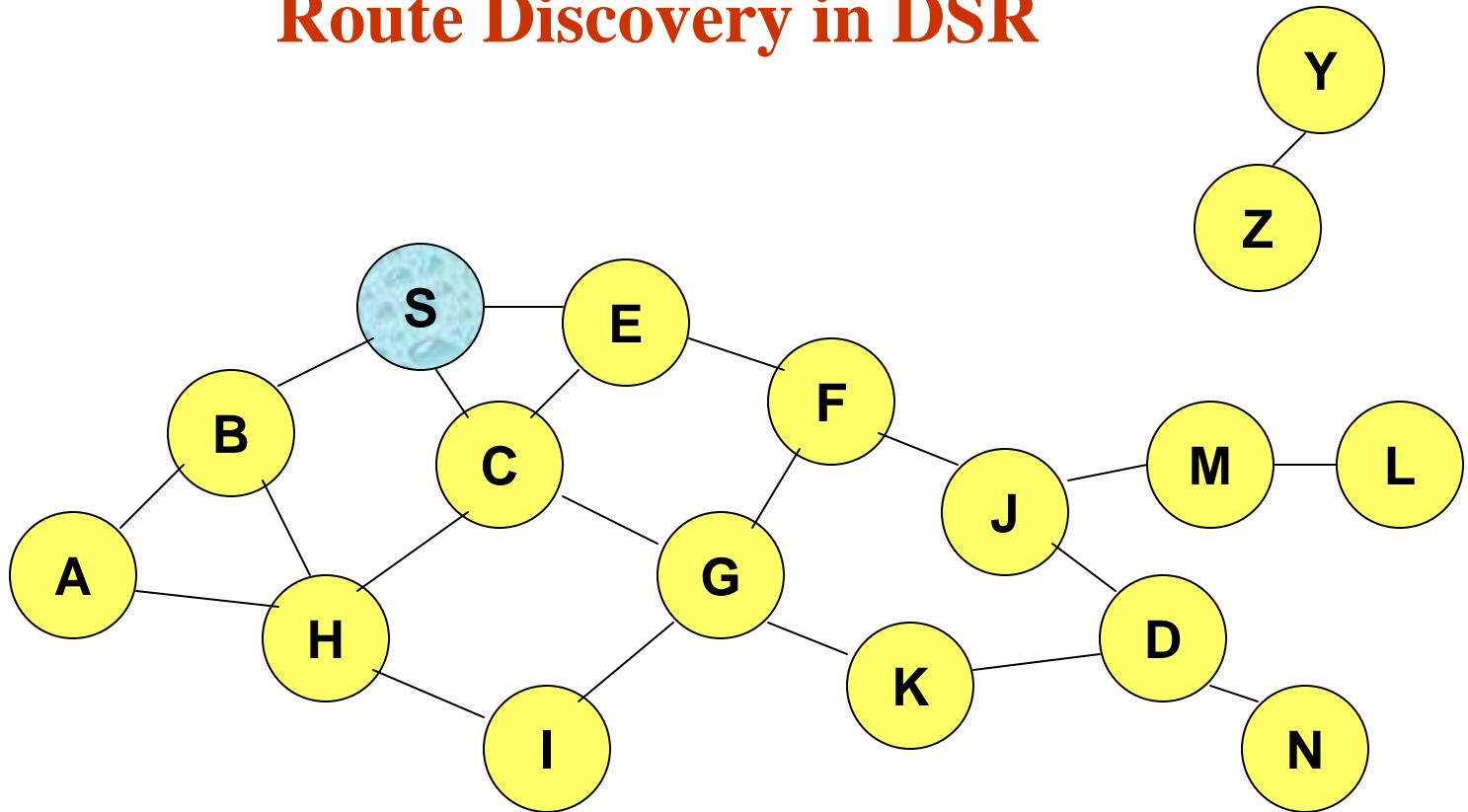
# Dynamic Source Routing (DSR) [Johnson96]

Courtesy of Prof. A. Abouz and N. Vaidya

## Dynamic Source Routing (DSR) [Johnson96]

- When node S wants to send a packet to node D, but does not know a route to D, node S initiates a **route discovery**
- Source node S floods **Route Request (RREQ)**
- Each node **appends own identifier** when forwarding RREQ

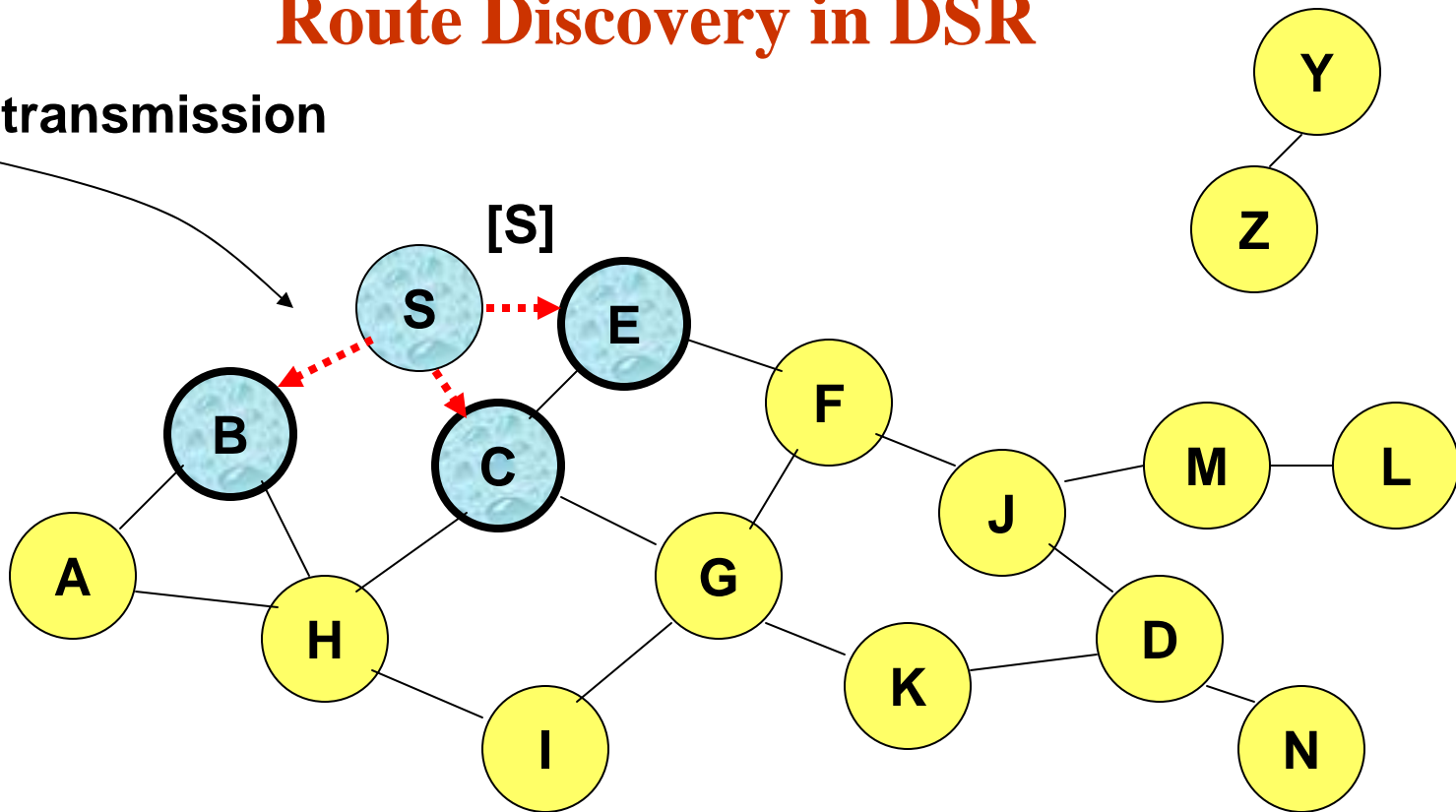
## Route Discovery in DSR



Represents a node that has received RREQ for D from S

# Route Discovery in DSR

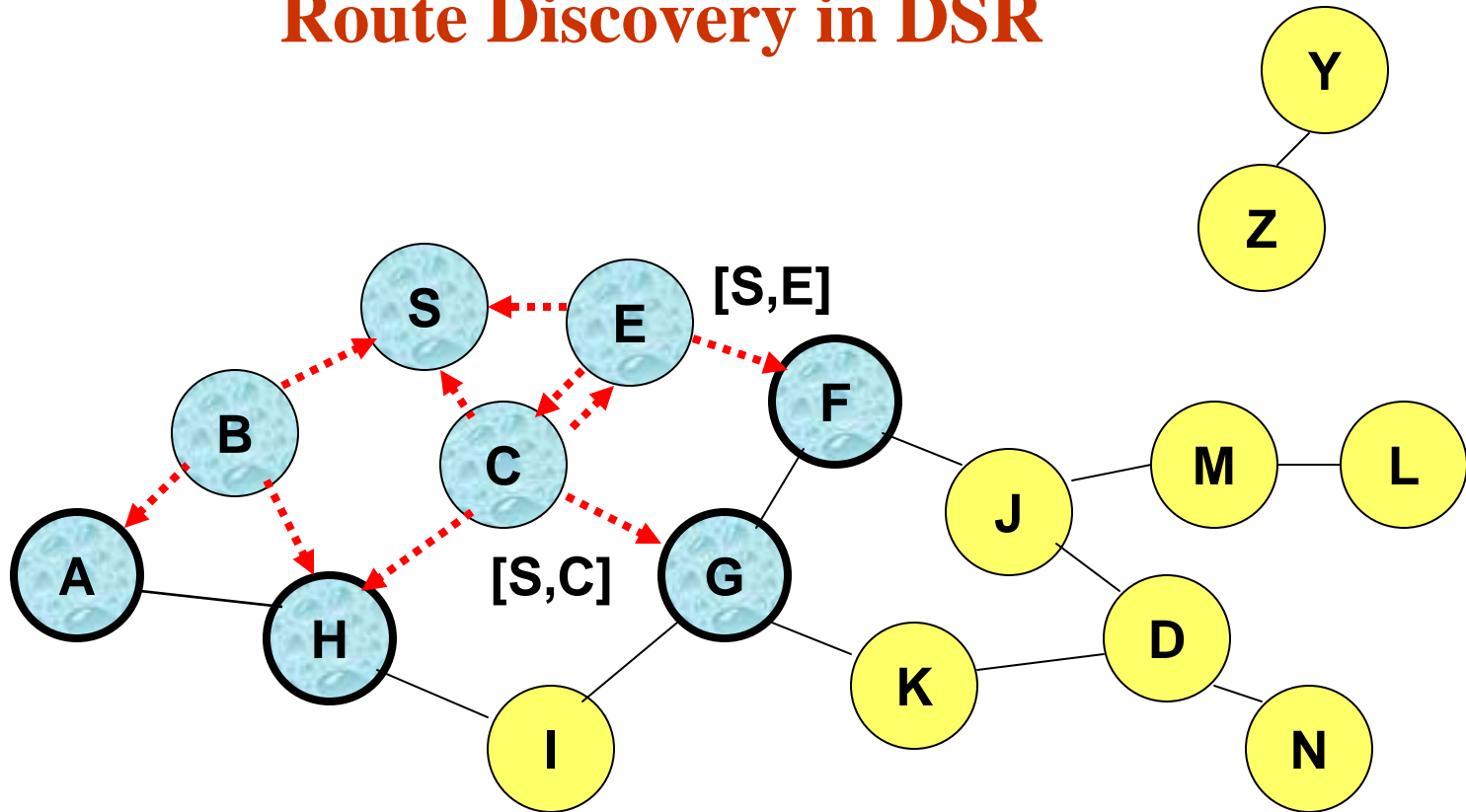
Broadcast transmission



.....→ Represents transmission of RREQ

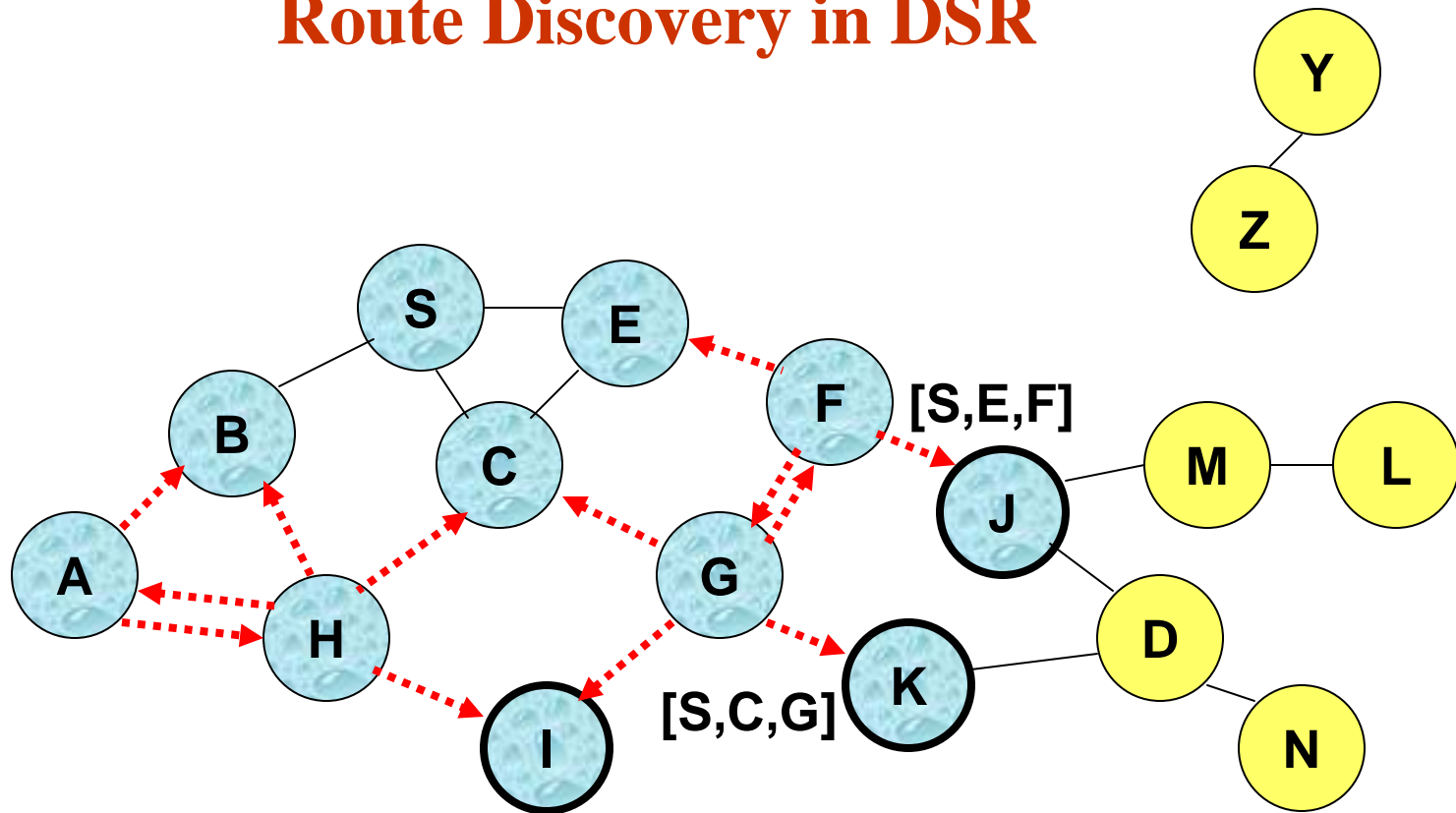
[X,Y] Represents list of identifiers appended to RREQ

## Route Discovery in DSR



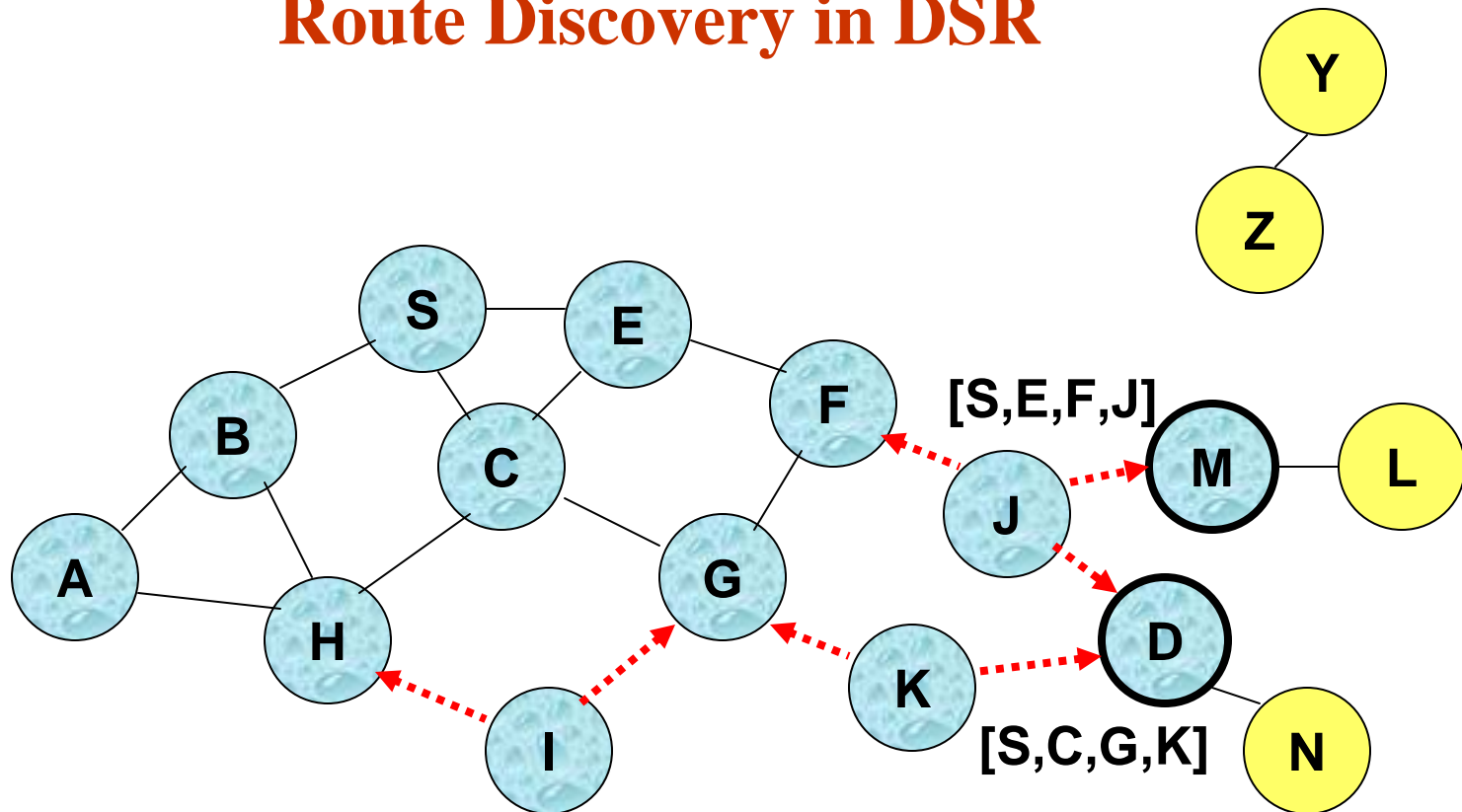
- Node H receives packet RREQ from two neighbors:  
**potential for collision**

## Route Discovery in DSR



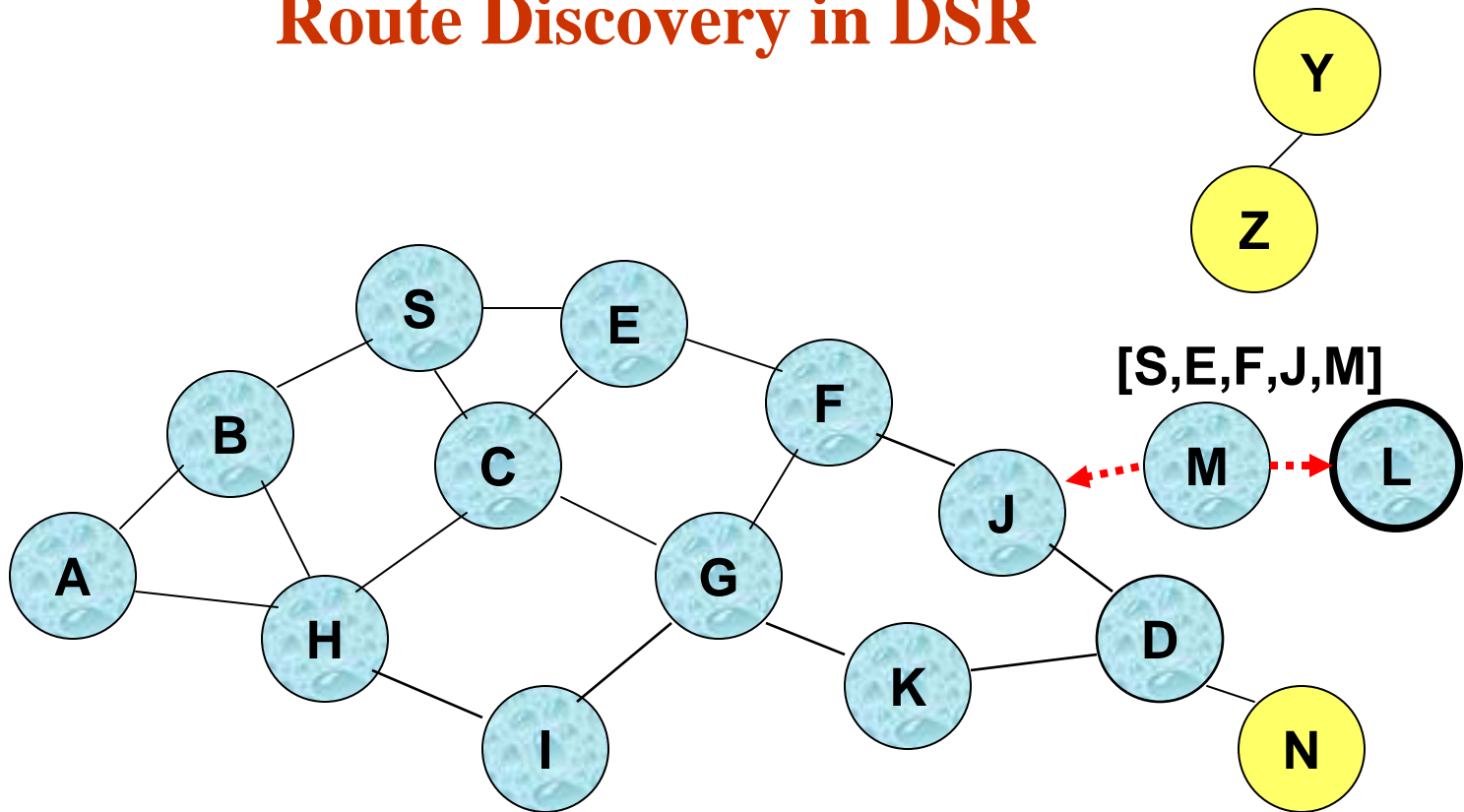
- Node C receives RREQ from G and H, but does not forward it again, because node C has **already forwarded RREQ** once

## Route Discovery in DSR



- Nodes J and K both broadcast RREQ to node D
- Since nodes J and K are **hidden** from each other, their **transmissions may collide**

## Route Discovery in DSR



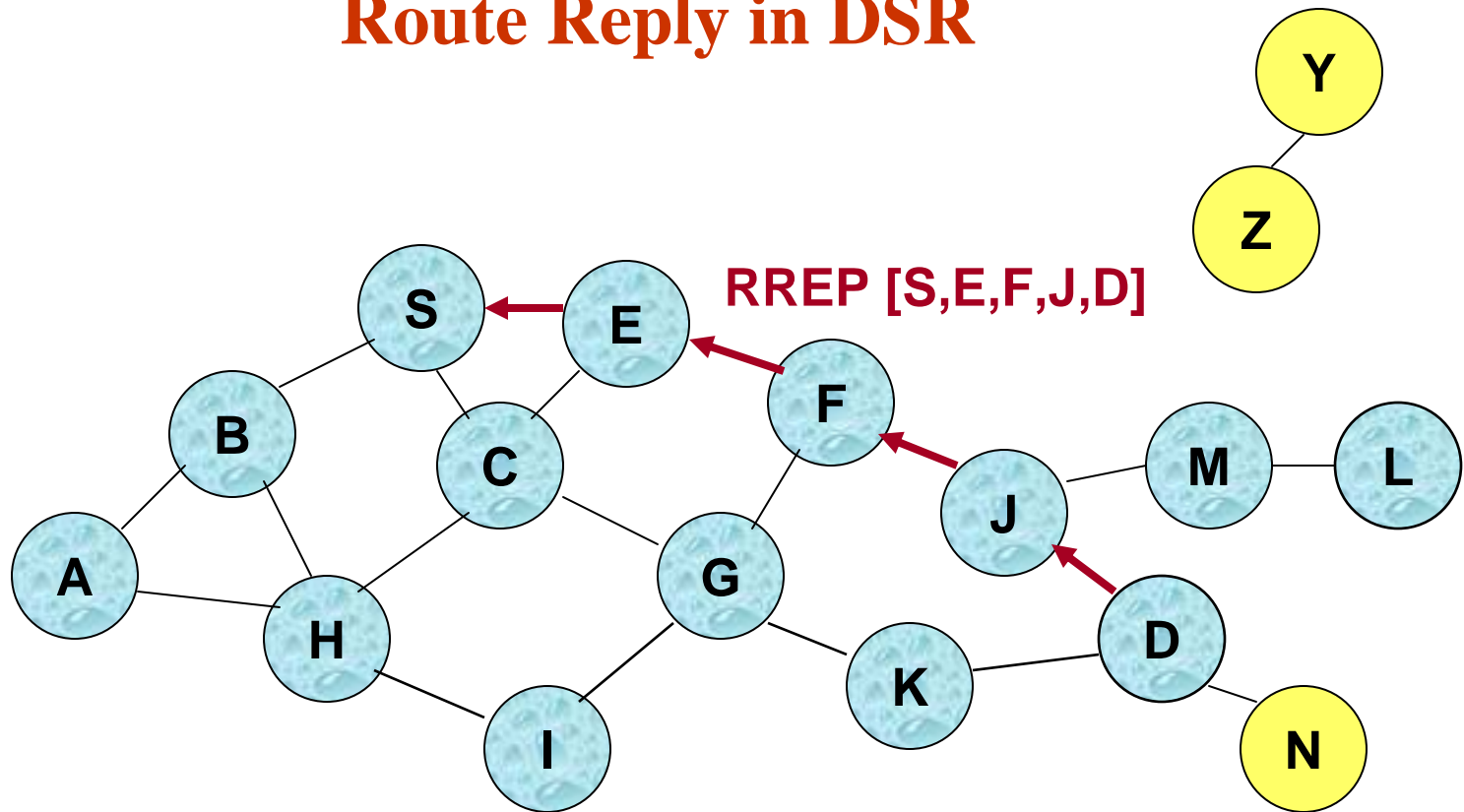
- Node D **does not forward** RREQ, because node D is the **intended target** of the route discovery



## Route Discovery in DSR

- Destination D on receiving the first RREQ, sends a **Route Reply (RREP)**
- RREP is sent on a route obtained by **reversing** the route appended to received RREQ
- RREP packet **contains the route** from S to D that was discovered using the RREQ packet

## Route Reply in DSR

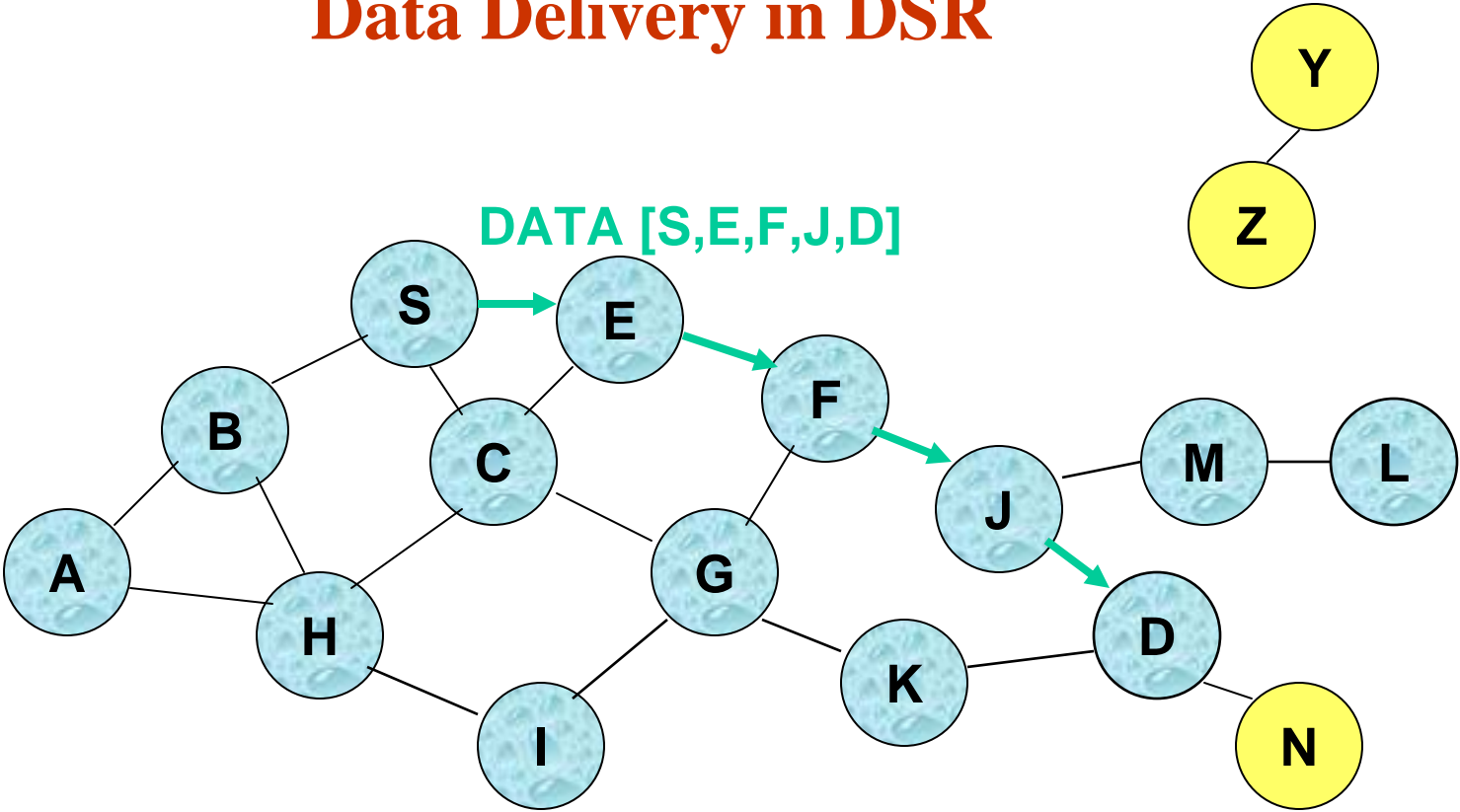


← Represents RREP control message

## Dynamic Source Routing (DSR)

- Node S on receiving RREP, **cached** the route included in the RREP
- When node S sends a data packet to D, the entire route is included in the packet header
  - hence the name **source routing**
- Intermediate nodes use the **source route** included in a packet to determine to whom a packet should be forwarded

# Data Delivery in DSR



Packet header size grows with route length

## When to Perform a Route Discovery?

- When node S wants to send data to node D (i.e. *on-demand*), but does not know a valid route to node D

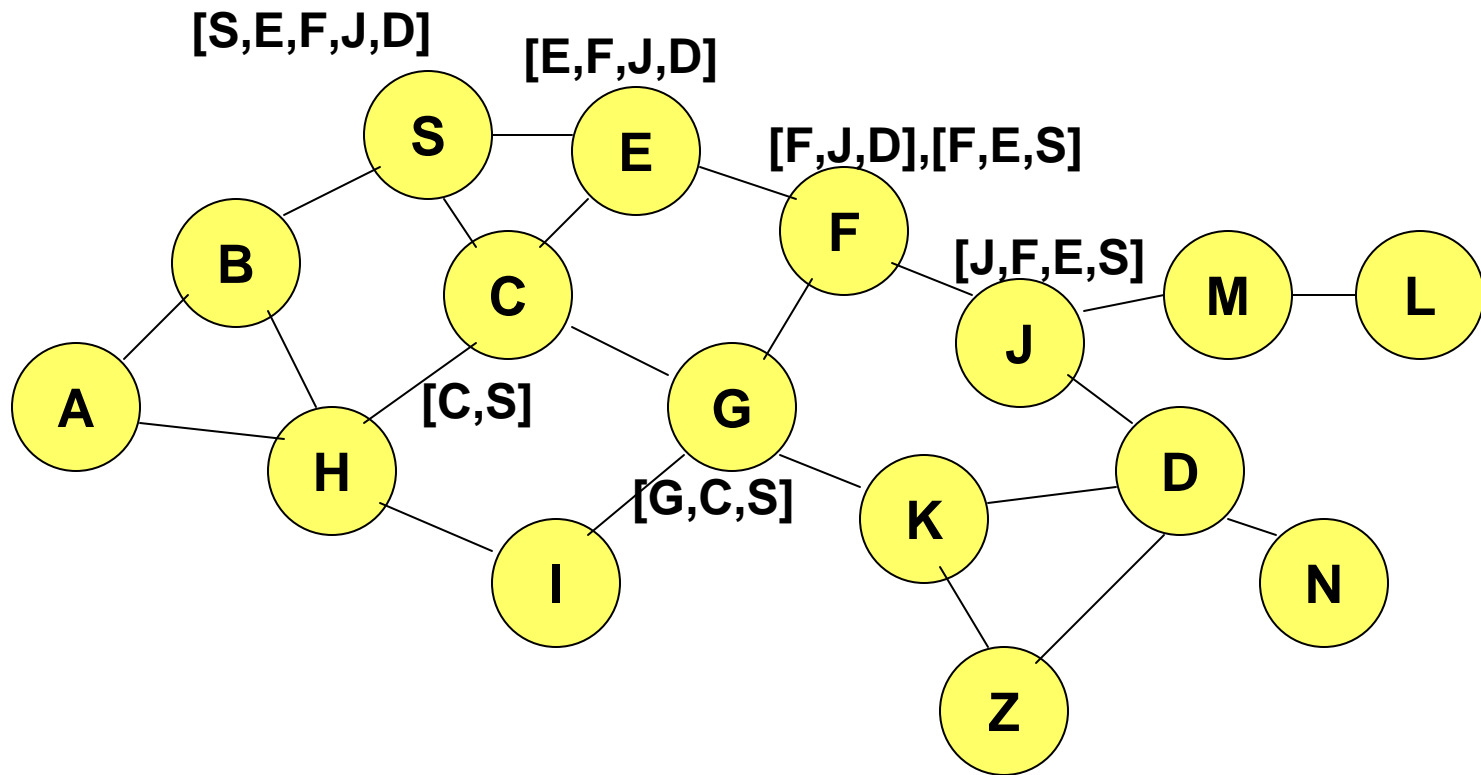
## DSR Optimization: Route Caching

- Each node caches a new route it learns by *any means*
- When node S finds route [S,E,F,J,D] to node D, node S also learns route [S,E,F] to node F
- When node K receives Route Request [S,C,G] destined for node D, node K learns route [K,G,C,S] to node S
- When node F forwards Route Reply [S,E,F,J,D], node F learns route [F,J,D] to node D
- When node E forwards Data [S,E,F,J,D] it learns route [E,F,J,D] to node D
- A node may also learn a route when it overhears Data packets!

## Use of Route Caching

- When node S learns that a route to node D is broken, it uses another route from its local cache, if such a route to D exists in its cache. Otherwise, node S initiates route discovery by sending a route request
- Node X on receiving a Route Request for some node D can send a Route Reply if node X knows a route to node D
- Use of route cache
  - can speed up route discovery
  - can reduce propagation of route requests

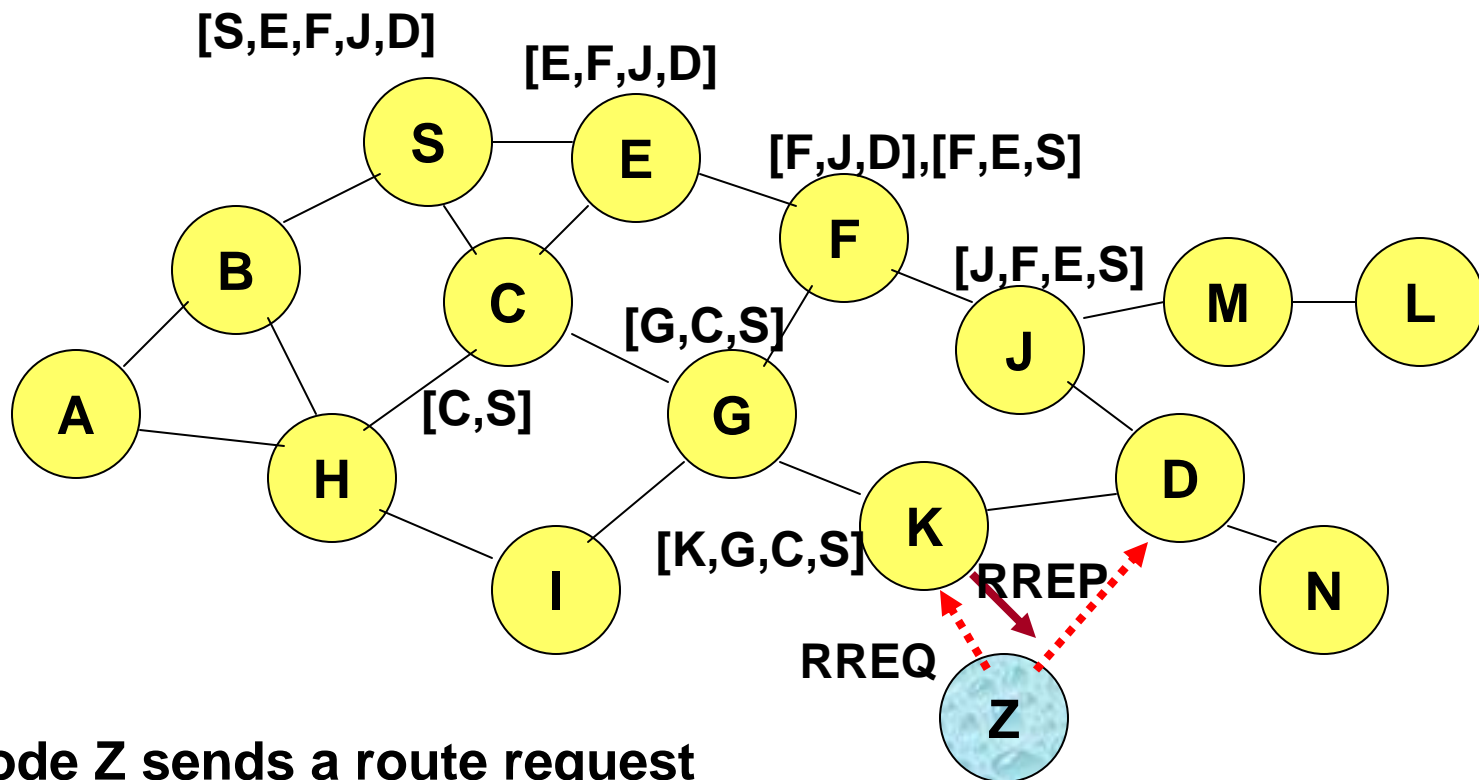
## Use of Route Caching



[P,Q,R] Represents cached route at a node

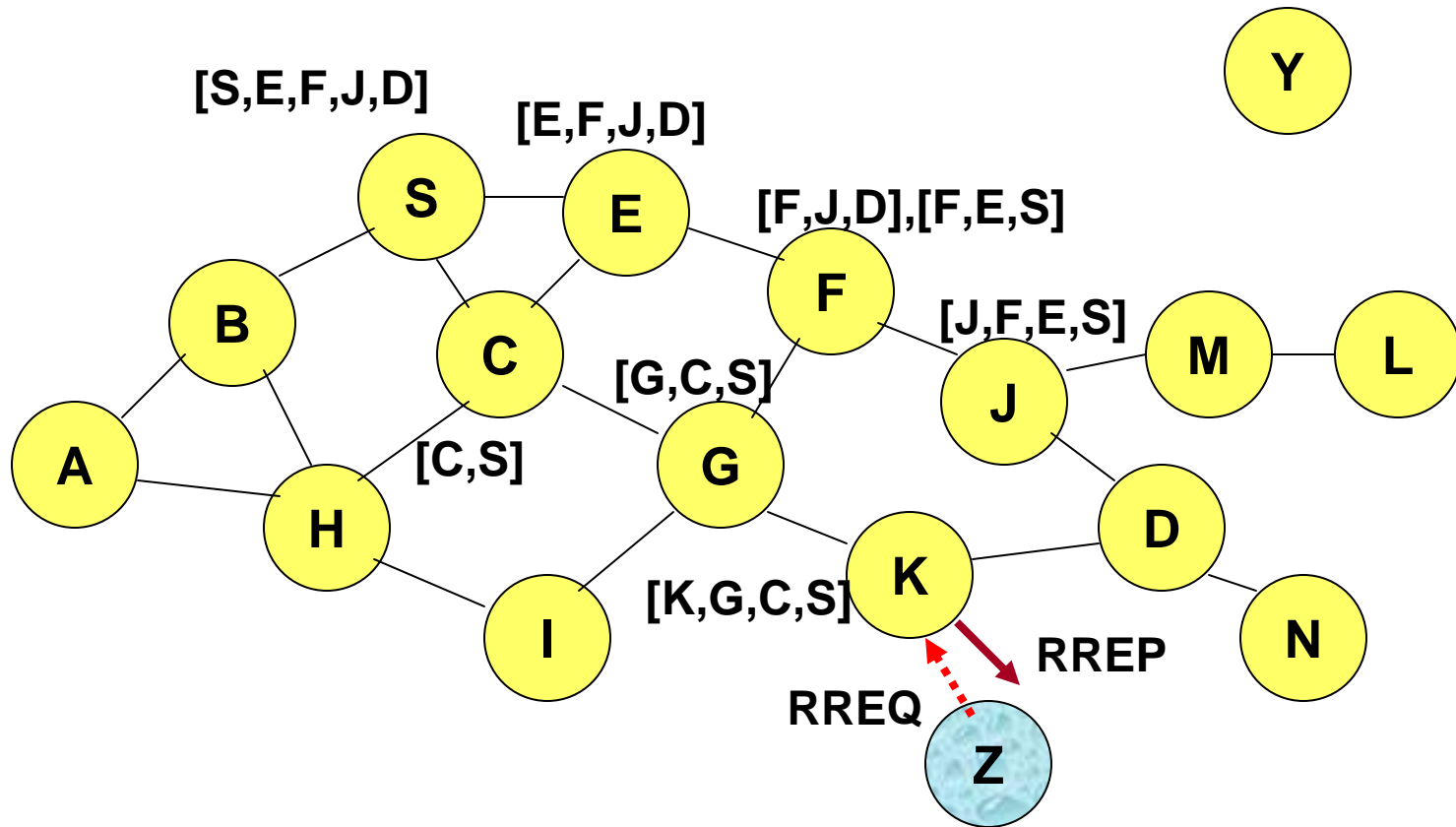


## Use of Route Caching: Can Speed up Route Discovery



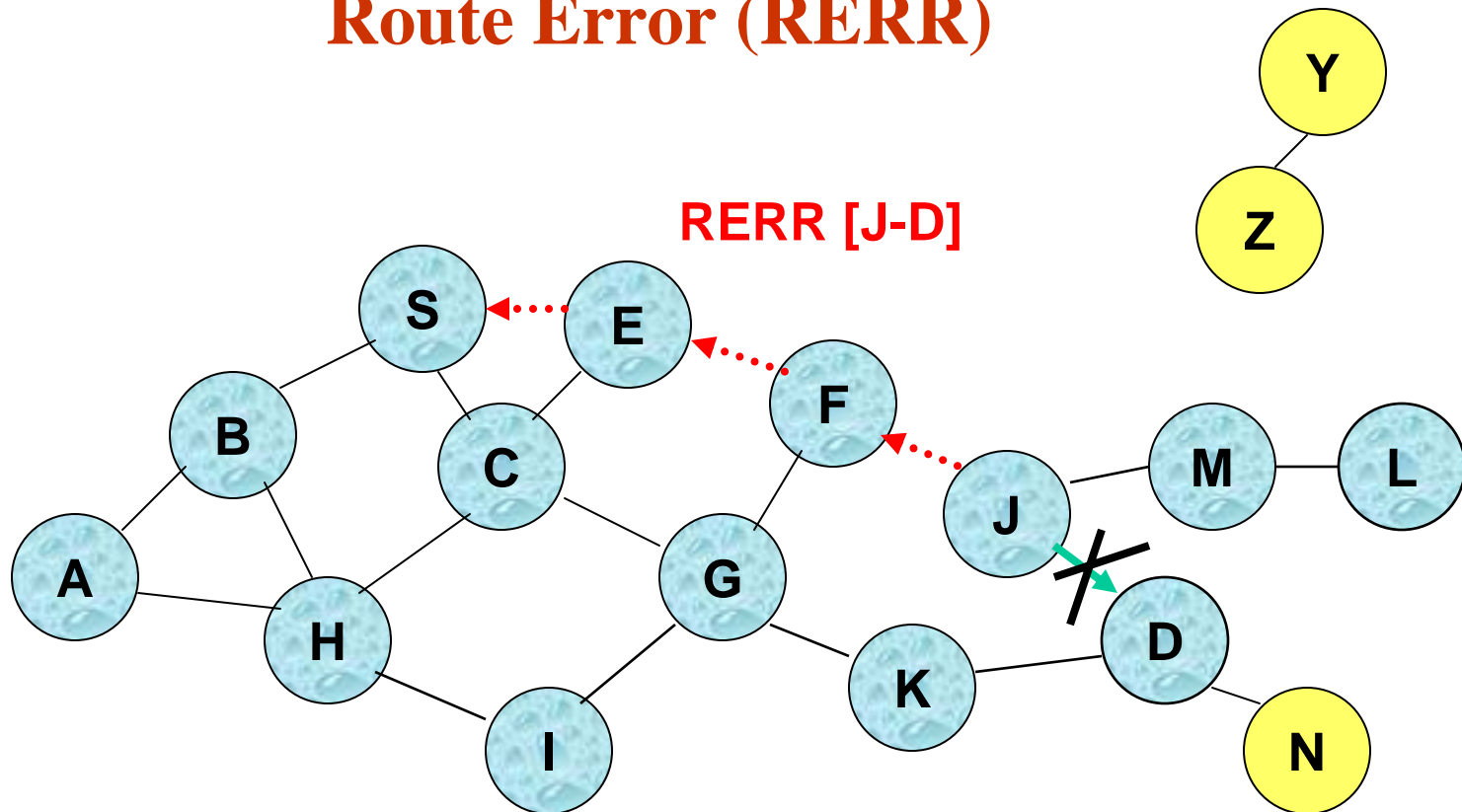
When node Z sends a route request for node C, node K sends back a route reply [Z,K,G,C] to node Z using a locally cached route

## Use of Route Caching: Can Reduce Propagation of Route Requests



Assume that there is no link between D and Z.  
Route Reply (RREP) from node K **limits flooding** of RREQ.  
In general, the reduction may be less dramatic.

## Route Error (RERR)



J sends a route error to S along route J-F-E-S when its attempt to forward the data packet S (with route SEFJD) on J-D fails

Nodes hearing RERR update their route cache to remove link J-D <sup>36</sup>

## Route Caching:Disadvantages

- Stale caches can adversely affect performance
- With passage of time and host mobility, cached routes may become invalid
- A sender host may try several stale routes (obtained from local cache, or replied from cache by other nodes), before finding a good route
- (An illustration of the adverse impact on TCP can be found in [\[Holland99\]](#))

## Dynamic Source Routing: Advantages

- Routes maintained only between nodes who need to communicate (ie. *on-demand*)
  - reduces overhead of route maintenance
- Route caching can further reduce route discovery overhead
- A single route discovery may yield many routes to the destination, due to multiple intermediate nodes replying from local caches

## Dynamic Source Routing: Disadvantages

- Packet header size grows with route length due to source routing
- Flood of route requests may potentially reach all nodes in the network
- Care must be taken to avoid collisions between route requests propagated by neighboring nodes
  - insertion of random delays before forwarding RREQ

## Dynamic Source Routing: Disadvantages (cache)

- Increased contention if too many route replies come back due to nodes replying using their local cache
  - *Route Reply Storm problem*
  - Reply storm may be eased by preventing a node from sending RREP if it hears another RREP with a shorter route
- An intermediate node may send Route Reply using a stale cached route, thus polluting other caches
- This problem can be eased if some mechanism to purge (potentially) invalid cached routes is incorporated.
- For some proposals for cache invalidation, see [\[Hu00Mobicom\]](#)
  - Static timeouts
  - Adaptive timeouts based on link stability