Abstract Data Types Documentation

Documentation

- Users are only interested in the properties of the ADT
- Programmers and designers require all the information which a user needs AND all information pertaining to the design and implementation
- Useful to think of the documentation as being an annotated definition of an abstract data type

Documentation Table of Contents

- Cover page, table of contents and abstract
- Document introduction
 - » Informal overview of the facilities provided. Help readers determine if this is what they need
- Data type objects
 - **»** Description of all the objects include diagrams
 - » Split into
 - > Imported which predefined objects are used
 - > Exported for others to use
 - > Hidden used in the implementation

- Operations
 - » Give
 - > Signature
 - > Informal description
 - > pre- and post- conditions
 - » Use natural language, mathematics, diagrams whatever best gets the meaning across.
 - » Be simple, complete, clear, precise, concise as possible

- Example partial axiomatic description of bank accounts
 - » The operation signatures only no pre- post- given

new : [] -> account

– Create an account with a zero balance

withdraw : account X amount -> account

- Remove amount from account
- deposit : account X amount -> account
 - Add amount to account

balance : account -> amount

– What is the amount in the account?

- Operation interaction
 - » Previous section describes operations in isolation
 - » Provide better understanding by showing properties when operations are used in combination
 - » Common descriptive method in use is axiomatic
 - > List of axioms or statements which must be true if the ADT is implemented and used correctly

- Axioms about the data type
 - » Axiom 1: New account has a balance of zero dollars balance(new) = 0
 - » Axiom 2: Cannot withdraw from a new account withdraw(new, amt) = error
 - » Axiom 3: Deposit amt and then withdraw amt with no intervening operations the balance does not change balance(withdraw(deposit(acct, amt) , amt) = balance(acct)
 - » Axiom 4: Only withdraw if the balance is ≥ the amount to withdraw. The amount is deducted from the balance balance(acct) < amt → withdraw(acct, amt) = error balance(acct) ≥ amt → balance(withdraw(acct, amt)) = balance(acct) amt</p>

- How to use the ADT
 - » Tutorial guide on use. Dwell on nuances. Describe various examples
- Dictionary
 - » Define new terminology or domain specific jargon that implementers or users may not know
- Undesired Event Dictionary
 - » Description of possible errors which can occur
 - » Contains warnings
 - » How to recognize error situations
 - **»** How to recover from error situations
 - » What to do if recovery is impossible

- ADT generation parameters
 - » Describe how instances and variations can be implemented from this generic data type
 - > How to change base types
 - > How to change amount of storage for a customer name
 - » Describe changes that can be made that will not violate assumptions and specifications. Design for a class of similar data types
 - » State what programming tools can be used to modify the implementation

- Design issues
 - » What were the design choices and why were the actual choices chosen. Help guide future changes to keep in the spirit of the original
 - > Why was fixed memory allocation used instead of dynamic?
 - > Why were size limits imposed?
 - > Why was a particular data structure chosen?

- Implementation notes
 - » Designer may have information of use to the implementer. Know properties that can improve implementation
- List of assumptions those assumptions that
 - » Cannot be violated
 - » Not implicit in the context
 - » Global
 - » Note: cannot state all assumptions so state those that
 - > Are most important
 - > Most likely to cause problems if violated
 - > Are not easily detected as causing problems until a long time later

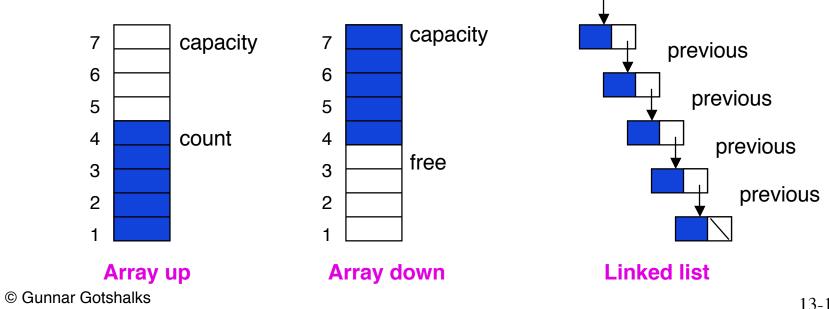
TCO - 10

- Normal use assumptions
 - » Information available from the ADT
 - » Information that must be supplied to the ADT
 - » Events reported by the ADT
 - » Tasks that can be performed by the ADT
 - » Operating states of the ADT and how they affect the Information obtained from and supplied to the ADT
 - » Failure states of the ADT and how they affect the information obtained from and supplied to the ADT

- Incorrect use assumptions
 - » Associated with run time undesired events
 - » What may or may not happen if the production version has undesired event handling code removed to speed up the system
- Program source text
 - » If the source test is small may be included with the description of the operations
- Facilities index
 - » A quick look up reference of all programs, modules, operations, objects and terms defined

Minimal Documentation

- Objects
 - » Types Diagrams where possible
- Example stack
 - » Imported none
 - » Exported STACK [G]
 - » Hidden implementation



last

Minimal Documentation – 2

- Operations example for a stack
 - >> Signatures, pre & post conditions > push : STACK [G] x G → STACK [G] - require true ensure result = x ^ s & count = old count + 1 > pop : STACK [G] → STACK [G] - require not empty (s) ensure result = s' & count = old count - 1 > top : STACK [G] → G - require not empty (s) ensure result = s₁

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- Operations example for a stack cont'd
 > empty : STACK [G] → BOOLEAN
 - require true
 ensure result = (count = 0)
 > new : [] → STACK [G]
 - require true ensure result = STACK [G] & count = 0
 - » Note: often "require true" is not written but is assumed
 - » It is better to write it as then one can wonder if it was left out by accident
 - > "nothing" is often represented with a special symbol. e.g. nil , λ , ϵ , Δ

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- Operations example for a stack cont'd
 - » axioms
 - > ∀ x : G, s : STACK [G] •

∧ empty (new)

top (push (s, x)) = x "it is the case that"
^ pop (push (s, x)) = s

 $\wedge \sim \text{empty}(\text{push}(s, x))$

is read as
 = X
 "it is the case."

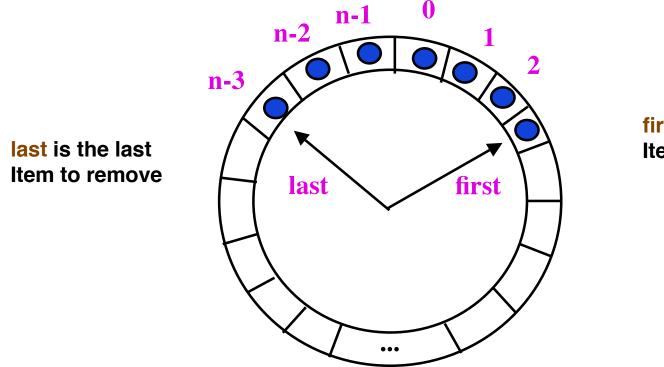
» Alternately can use natural language > forall x : G, s : STACK [G } :: top (push (s, x)) = x and pop (push (s, x)) = s and empty (new) and not empty (push (s, x))

ADT Invariants

- Conditions that must be true after the execution of any method in the the class
- The conditions that hold, at all times, among the objects in an instance of the ADT
 - » More on this when we discuss design by contract

Example Circular Queue

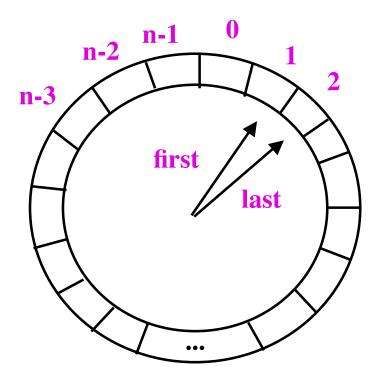
isEmpty \rightarrow length = 0 & (last-1) mod Size = first isFull \rightarrow length = Size - 1 not isFull \rightarrow length = (Size + first - last + 1) mod Size



first is the first Item to remove

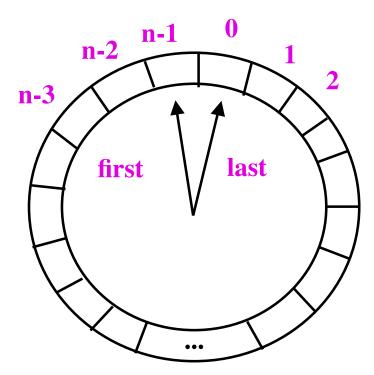
Empty Circular Queue

isEmpty → length = 0 & (last-1) mod Size = first isFull → length = Size - 1 not isFull → length = (Size + first - last + 1) mod Size



Empty Circular Queue – 2

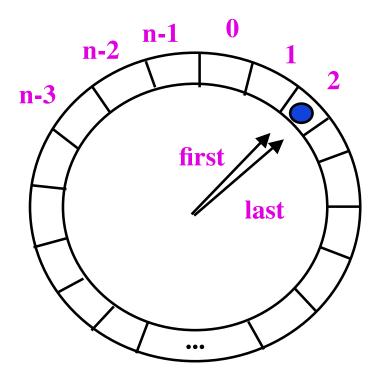
isEmpty \rightarrow length = 0 & (last-1) mod Size = first isFull \rightarrow length = Size - 1 not isFull \rightarrow length = (Size + first - last + 1) mod Size



length = $(n + (n-1) - 0 + 1) \mod n$ = $(2n + 0) \mod n$ = 0

Length 1 Circular Queue

isEmpty \rightarrow length = 0 & (last-1) mod Size = first isFull \rightarrow length = Size - 1 not isFull \rightarrow length = (Size + first - last + 1) mod Size



Longer length Circular Queue

length = (Size + first - last + 1) mod Size

