

## Homework Assignment #2

### Due: October 5, 2010

1. Recall the formal definition of a string from class: a string over alphabet  $\Sigma$  is a function from  $\{i \in \mathbb{N} : 1 \leq i \leq \ell\}$  to  $\Sigma$  (where  $\ell$  is some natural number called the length of the string).

Recall the formal definition of string concatenation from class: If  $s_1$  and  $s_2$  are strings over  $\Sigma$  of length  $\ell_1$  and  $\ell_2$ , respectively, then  $s = s_1 \cdot s_2$  is a string over  $\Sigma$  of length  $\ell_1 + \ell_2$  defined as follows.

$$s(i) = \left\{ \begin{array}{ll} s_1(i) & \text{if } 1 \leq i \leq \ell_1 \\ s_2(i - \ell_1) & \text{if } \ell_1 + 1 \leq i \leq \ell_1 + \ell_2. \end{array} \right\}$$

If  $z$  is a string of length  $\ell$  over the alphabet  $\Sigma$ , then the reverse of  $z$ , denoted  $z^R$ , is a string of length  $\ell$  over  $\Sigma$  defined as follows.

$$z^R(i) = z(\ell - i + 1).$$

Use these definitions to give a careful proof that, for any strings  $x$  and  $y$  over  $\Sigma$ ,  $(x \cdot y)^R = y^R \cdot x^R$ .

2. Draw a deterministic finite automaton that decides the language of all binary strings that do not contain 1010 as a substring. You do not have to prove your answer is correct. However, for each state of your automaton, write down a description of all strings that take the machine to that state.

**Optional programming task** (Do *not* hand this in.) Write a Java programme which, given a description of a deterministic finite automaton and an input string, determines whether the automaton accepts or rejects the string. Assume that the automaton description is given in the following format.

- The first line gives two integers separated by a space:  $n$ , the number of states, and  $m$ , the number of accepting states. (We will assume the states are labelled  $1, 2, \dots, n$  and that state 1 is the starting state.)
- The second line gives a string containing one copy of each character in the alphabet.
- The third line gives a list of the  $m$  accepting states, separated by spaces. (If there are no accepting states, this line will be blank.)
- The remaining lines describe the possible transitions. If  $\delta(q, a) = q'$ , there will be a line containing  $q, a$  and  $q'$ , separated by spaces. The end of the transitions will be indicated by a line containing just 0.

For example, the automaton shown on page 34 of the textbook would be specified as follows:

```
3 1
01
2
1 0 1
1 1 2
2 0 3
2 1 2
3 0 2
3 1 2
0
```